

# Hood-Melville Queue

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## Abstract

This is a verified implementation of a constant time queue. The original design is due to Hood and Melville [1]. This formalization follows the presentation by Okasaki [2].

```
theory Hood-Melville-Queue
imports
  HOL-Data-Structures.Queue-Spec
begin

datatype 'a status =
  Idle
  | Rev nat 'a list 'a list 'a list 'a list
  | App nat 'a list 'a list
  | Done 'a list

record 'a queue = lenf :: nat
  front :: 'a list
  status :: 'a status
  rear :: 'a list
  lenr :: nat

fun exec :: 'a status  $\Rightarrow$  'a status where
  exec (Rev ok (x#f) f' (y#r) r') = Rev (ok+1) f (x#f') r (y#r')
| exec (Rev ok [] f' [y] r') = App ok f' (y#r')
| exec (App 0 f' r') = Done r'
| exec (App ok (x#f') r') = App (ok-1) f' (x#r')
| exec s = s

fun invalidate where
```

```

| invalidate (Rev ok f f' r r') = Rev (ok-1) f f' r r'
| invalidate (App 0 f' (x#r')) = Done r'
| invalidate (App ok f' r') = App (ok-1) f' r'
| invalidate s = s

```

**fun** *exec2* :: 'a queue ⇒ 'a queue **where**

```

exec2 q =
  (case exec (exec (status q)) of
    Done newf ⇒ q(status := Idle, front := newf) |
    newstate ⇒ q(status := newstate))

```

**definition** *check* :: 'a queue ⇒ 'a queue **where**

```

check q = (if lenr q ≤ lenf q
  then exec2 q
  else let newstate = Rev 0 (front q) [] (rear q) []
    in exec2 (q(lenf := lenf q + lenr q, status := newstate, rear := [],
lenr := 0)))

```

**definition** *empty* :: 'a queue **where**

```

empty = queue.make 0 [] Idle [] 0

```

**fun** *enq* **where**

```

enq x q = check (q(rear := x#(rear q), lenr := lenr q + 1))

```

**fun** *deq* **where**

```

deq q = check (q(lenf := lenf q - 1
, front := tl (front q)
, status := invalidate (status q)))

```

**fun** *front-list* :: 'a queue ⇒ 'a list **where**

```

front-list q = (case status q of
  Idle ⇒ front q
| Done f ⇒ f
| Rev ok f f' r r' ⇒ rev (take ok f') @ f @ rev r @ r'
| App ok f' r' ⇒ rev (take ok f') @ r')

```

**definition** *rear-list* :: 'a queue ⇒ 'a list **where**

```

rear-list = rev o rear

```

**fun** *list* :: 'a queue ⇒ 'a list **where**

```

list q = front-list q @ rear-list q

```

**fun** *first* :: 'a queue  $\Rightarrow$  'a **where**  
*first* q = hd (front q)

**fun** *rem-steps* :: 'a status  $\Rightarrow$  nat **where**  
*rem-steps* (Rev ok f f' r r') = 2\*length f + ok + 2  
| *rem-steps* (App ok f' r') = ok + 1  
| *rem-steps* - = 0

**fun** *inv-st* :: 'a status  $\Rightarrow$  bool **where**  
*inv-st* (Rev ok f f' r r') = (length f + 1 = length r  $\wedge$   
length f' = length r'  $\wedge$   
ok  $\leq$  length f')  
| *inv-st* (App ok f' r') = (ok  $\leq$  length f'  $\wedge$  length f' < length r')  
| *inv-st* - = True

**fun** *steps* :: nat  $\Rightarrow$  'a status  $\Rightarrow$  'a status **where**  
*steps* n st = (exec  $\overset{\sim}{\sim}$  n) st

**lemma** *rev-steps-app*:

**assumes** *inv*: *inv-st* (Rev ok f f' r r')  
**shows** *steps* (length f + 1) (Rev ok f f' r r') = App (length f + ok) (rev f @ f')  
(rev r @ r')  
<proof>

**lemma** *inv-st-steps*:

**assumes** *inv* : *inv-st* s  
**assumes** *not-idle* : s  $\neq$  Idle  
**shows**  $\exists x. \text{steps } (\text{rem-steps } s) s = \text{Done } x$  (**is** ?reach-done s)  
<proof>

**lemma** *inv-st-exec*:

**assumes** *inv-st*: *inv-st* s  
**shows** *inv-st* (exec s)  
<proof>

**lemma** *inv-st-exec2*:

**assumes** *inv-st*: *inv-st* s  
**shows** *inv-st* (exec (exec s))  
<proof>

**lemma** *inv-st-invalidate*:

**assumes** *inv-st*: *inv-st s*  
**shows** *inv-st* (*invalidate s*)  
⟨*proof*⟩

**definition** *invar* **where**

*invar*  $q = (\text{lenf } q = \text{length } (\text{front-list } q) \wedge$   
 $\text{lenr } q = \text{length } (\text{rear-list } q) \wedge$   
 $\text{lenr } q \leq \text{lenf } q \wedge$   
*(case status*  $q$  *of*  
 $\text{Rev } ok \text{ } f \text{ } f' \text{ } r \text{ } r' \Rightarrow 2 * \text{lenr } q \leq \text{length } f' \wedge ok \neq 0 \wedge 2 * \text{length } f + ok$   
 $+ 2 \leq 2 * \text{length } (\text{front } q)$   
 $| \text{App } ok \text{ } f \text{ } r \Rightarrow 2 * \text{lenr } q \leq \text{length } r \wedge ok + 1 \leq 2 * \text{length } (\text{front } q)$   
 $| - \Rightarrow \text{True}) \wedge$   
 $(\exists \text{rest. front-list } q = \text{front } q @ \text{rest}) \wedge$   
 $(\neg(\exists \text{fr. status } q = \text{Done } \text{fr})) \wedge$   
 $\text{inv-st } (\text{status } q))$

**lemma** *invar-empty*: *invar empty*  
⟨*proof*⟩

**lemma** *tl-rev-take*:  $\llbracket 0 < ok; ok \leq \text{length } f \rrbracket \Longrightarrow \text{rev } (\text{take } ok \text{ } (x \# f)) = \text{tl } (\text{rev } (\text{take } ok \text{ } f)) @ [x]$   
⟨*proof*⟩

**lemma** *tl-rev-take-Suc*:  
 $n + 1 \leq \text{length } l \Longrightarrow \text{rev } (\text{take } n \text{ } l) = \text{tl } (\text{rev } (\text{take } (\text{Suc } n) \text{ } l))$   
⟨*proof*⟩

**lemma** *invar-deq*:  
**assumes** *inv*: *invar*  $q$   
**shows** *invar* (*deq*  $q$ )  
⟨*proof*⟩

**lemma** *invar-enq*:  
**assumes** *inv*: *invar*  $q$   
**shows** *invar* (*enq*  $x$   $q$ )  
⟨*proof*⟩

**lemma** *queue-correct-deq* :  
**assumes** *inv*: *invar*  $q$   
**shows** *list* (*deq*  $q$ ) = *tl* (*list*  $q$ )  
⟨*proof*⟩

**lemma** *queue-correct-enq* :  
**assumes** *inv*: *invar* *q*  
**shows** *list* (*enq* *x* *q*) = (*list* *q*) @ [*x*]  
⟨*proof*⟩

**datatype** *'a* *action* = *Deq* | *Enq* *'a*

**type-synonym** *'a* *actions* = *'a* *action* *list*

**fun** *do-act* :: *'a* *action* ⇒ *'a* *queue* ⇒ *'a* *queue* **where**  
*do-act* *Deq* *q* = *deq* *q*  
| *do-act* (*Enq* *x*) *q* = *enq* *x* *q*

**definition** *qfa* :: *'a* *actions* ⇒ *'a* *queue* **where**  
*qfa* = (λ*acts*. *foldr* *do-act* *acts* *empty*)

**lemma** *invar-qfa* : *invar* (*qfa* *l*)  
⟨*proof*⟩

**lemma** *qfa-deq-correct*: *list* (*deq* (*qfa* *l*)) = *tl* (*list* (*qfa* *l*))  
⟨*proof*⟩

**lemma** *qfa-enq-correct*: *list* (*enq* *x* (*qfa* *l*)) = (*list* (*qfa* *l*)) @ [*x*]  
⟨*proof*⟩

**lemma** *first-correct* :  
**assumes** *inv*: *invar* *q*  
**assumes** *not-nil* : *list* *q* ≠ []  
**shows** *first* *q* = *hd* (*list* *q*)  
⟨*proof*⟩

**fun** *is-empty* :: *'a* *queue* ⇒ *bool* **where**  
*is-empty* *q* = (*list* *q* = [])

**interpretation** *HMQ*: *Queue* **where**  
*empty* = *empty* **and**  
*enq* = *enq* **and**  
*first* = *first* **and**  
*deq* = *deq* **and**  
*is-empty* = *is-empty* **and**  
*list* = *list* **and**  
*invar* = *invar*

*<proof>*

**end**

## **References**

- [1] R. Hood and R. Melville. Real-time queue operation in pure LISP. *Inf. Process. Lett.*, 13(2):50–54, 1981.
- [2] C. Okasaki. *Purely Functional Data Structures*. Cambridge University Press, 1998.