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# Llumnix: Dynamic Scheduling for Large Language Model Serving

Biao Sun\*, Ziming Huang\*, Hanyu Zhao\*, Wencong Xiao, Xinyi Zhang, Yong Li, Wei Lin (\* Equal contribution)





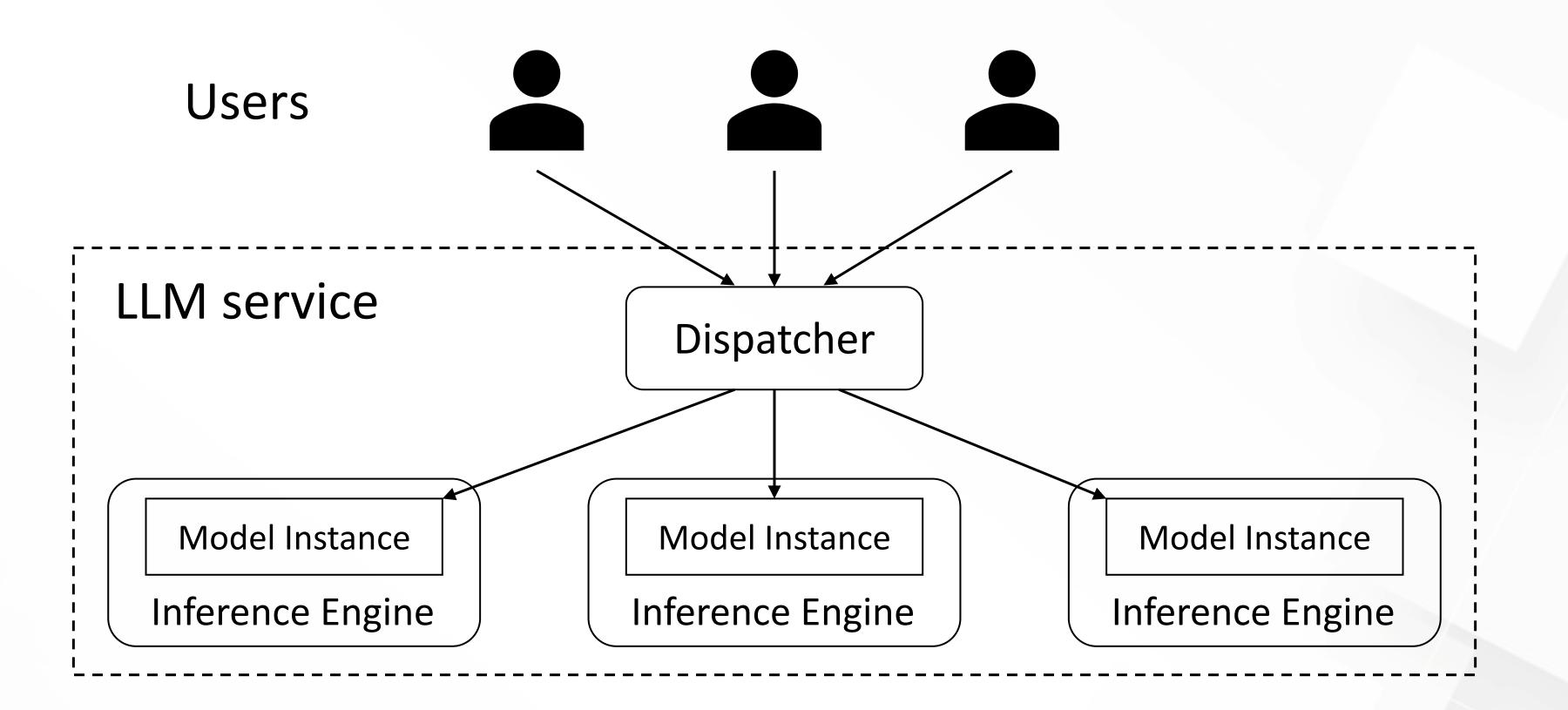
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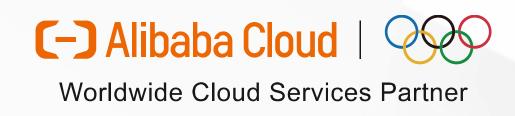


# **LLM Serving Today: A Cluster Perspective**

• A request dispatcher + multiple instances of an inference engine

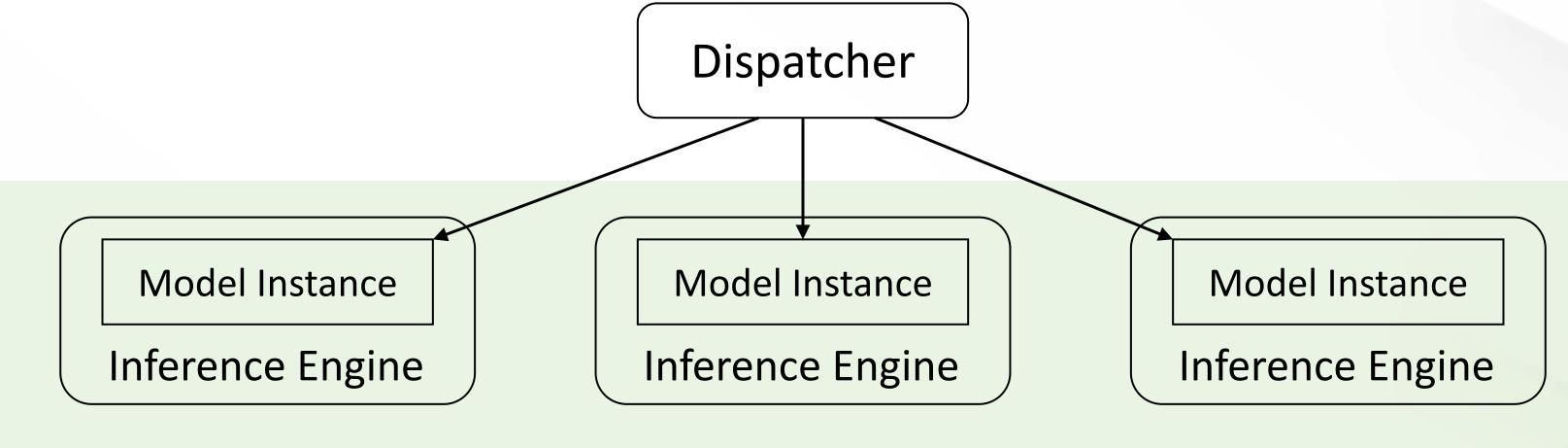






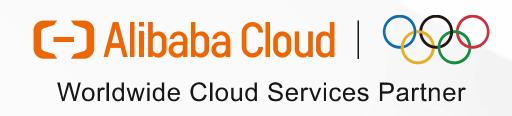
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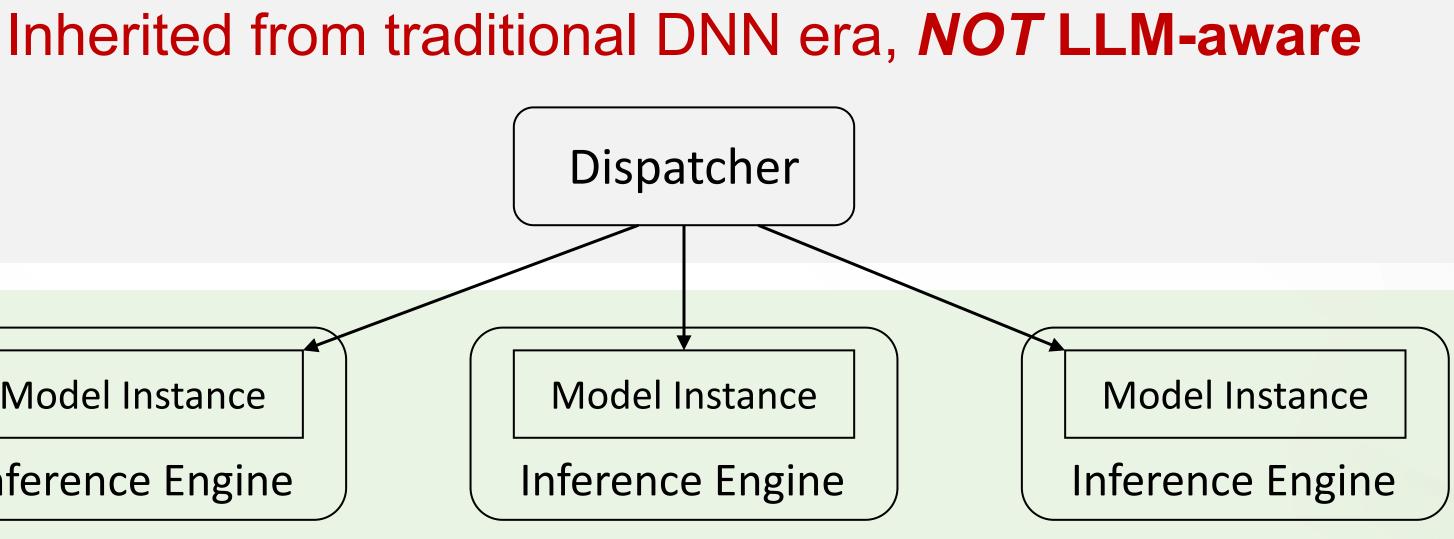
LLM-tailored, optimizing single-instance performance

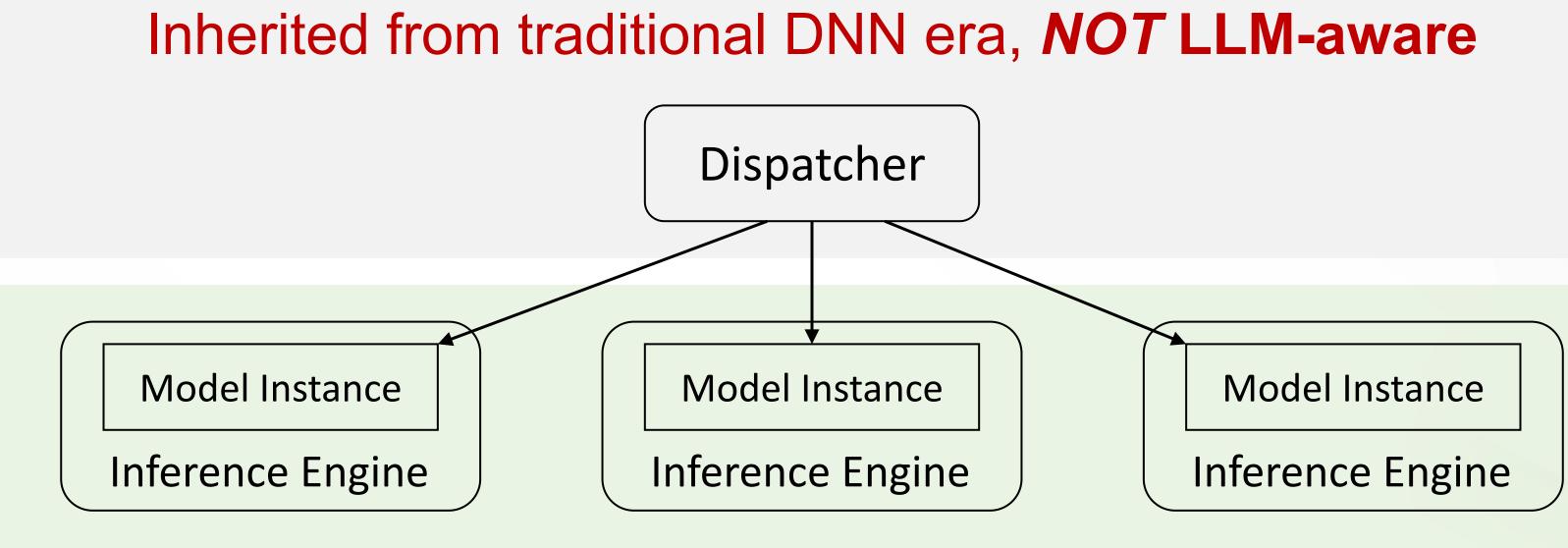




# LLM Serving Today: A Cluster Perspective

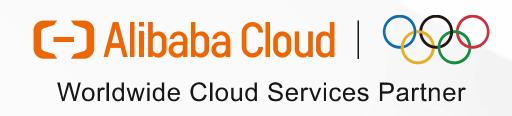
A request dispatcher + multiple instances of an inference engine





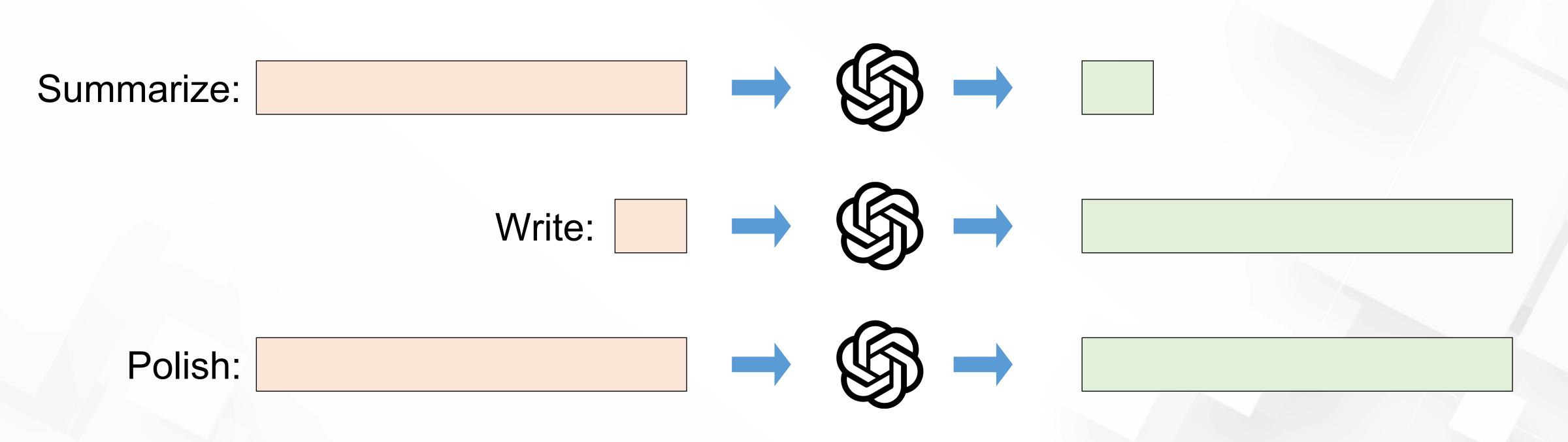
LLM-tailored, optimizing single-instance performance





# LLM Characteristic (1): Workload Heterogeneity

- Universal models, diverse applications
- Requests are heterogeneous ullet
  - Sequence (input/output) lengths ullet







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# LLM Characteristic (1): Workload Heterogeneity

- Universal models, diverse applications
- Requests are heterogeneous
  - Sequence (input/output) lengths ullet
  - Latency SLOs: interactive vs. offline, ChatGPT plus vs. normal •

# Introducing ChatGPT Plus

The new subscription plan, ChatGPT Plus, will be available for \$20/month, and subscribers will receive a number of benefits:

- General access to ChatGPT, even during peak times
- Faster response times



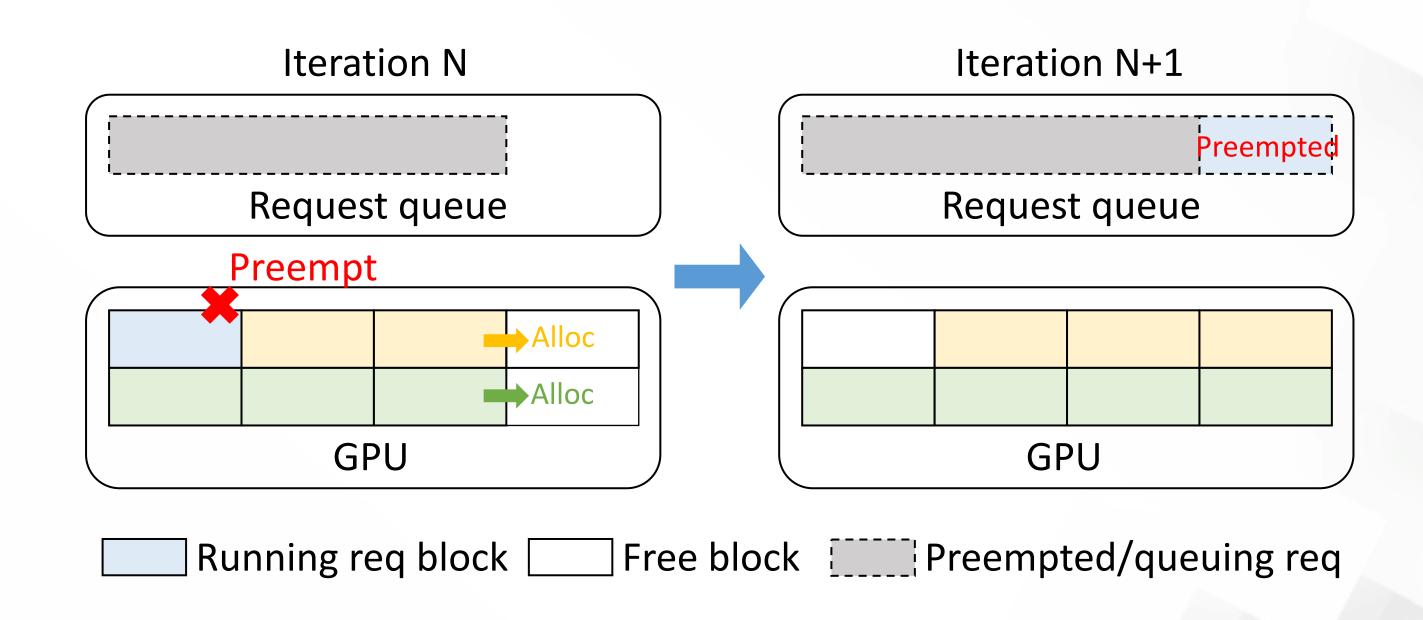
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# LLM Characteristic (2): *Execution Unpredictability*

- Autoregressive execution
  - Output lengths not known a priori •
  - Dynamic GPU memory demands of KV caches •



[1] Kwon et al. Efficient Memory Management for Large Language Model Serving with PagedAttention (SOSP '23)



# • State of the art: paged memory allocation + preemptive scheduling <sup>[1]</sup>

# **Challenge (1): Performance Isolation**

Preemptions -> poor tail latencies

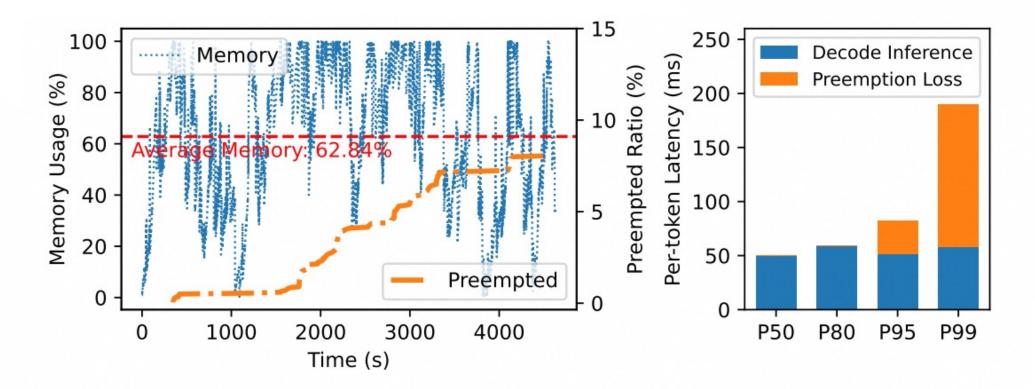
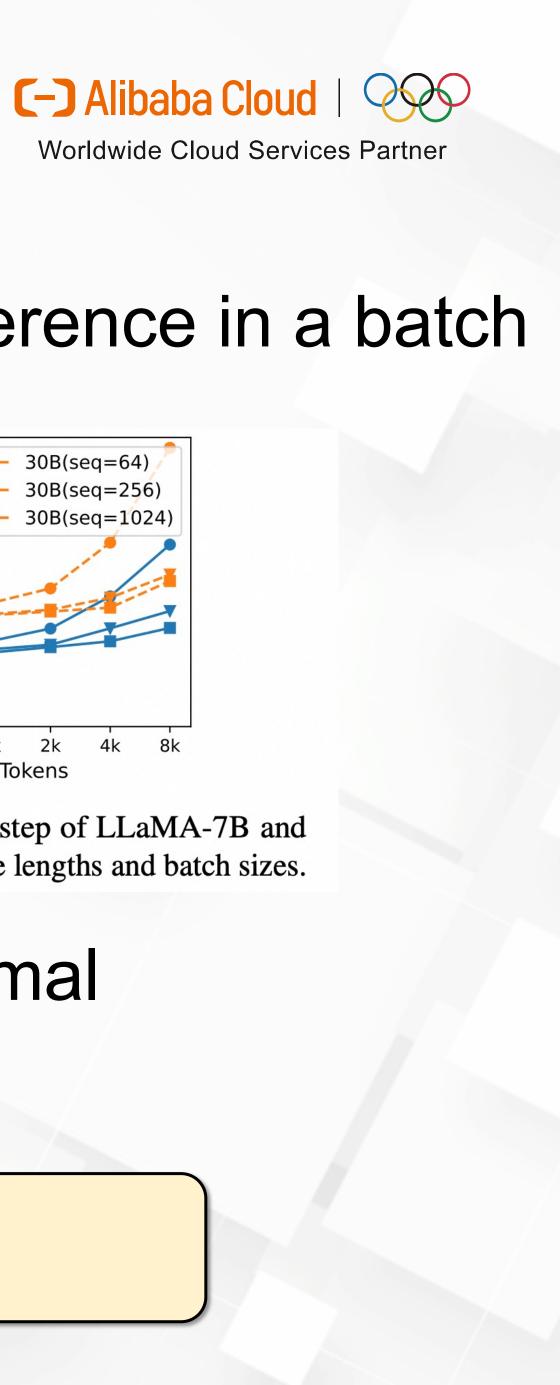


Figure 3: Request preemptions in LLaMA-7B serving.

 Load balancing via one-shot dispatching could be suboptimal due to unpredictable execution



## Performance interference in a batch

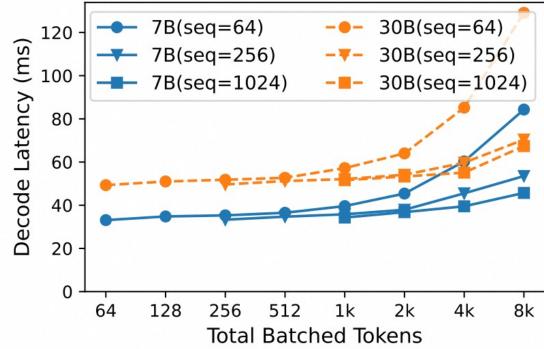
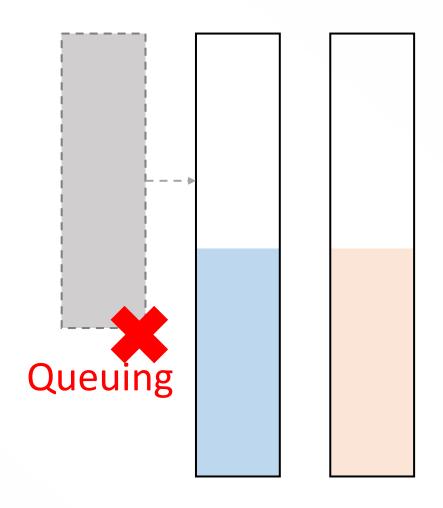


Figure 4: Latencies of one decode step of LLaMA-7B and LLaMA-30B with different sequence lengths and batch sizes.

## Requirement (1): Continuous load balancing

# Challenge (2): Memory Fragmentation

- Load balancing -> fragmentation across instances
  - A classic spreading vs. packing tradeoff

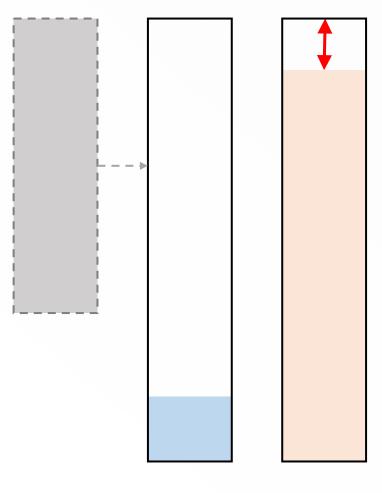


Spreading



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## across instances deoff

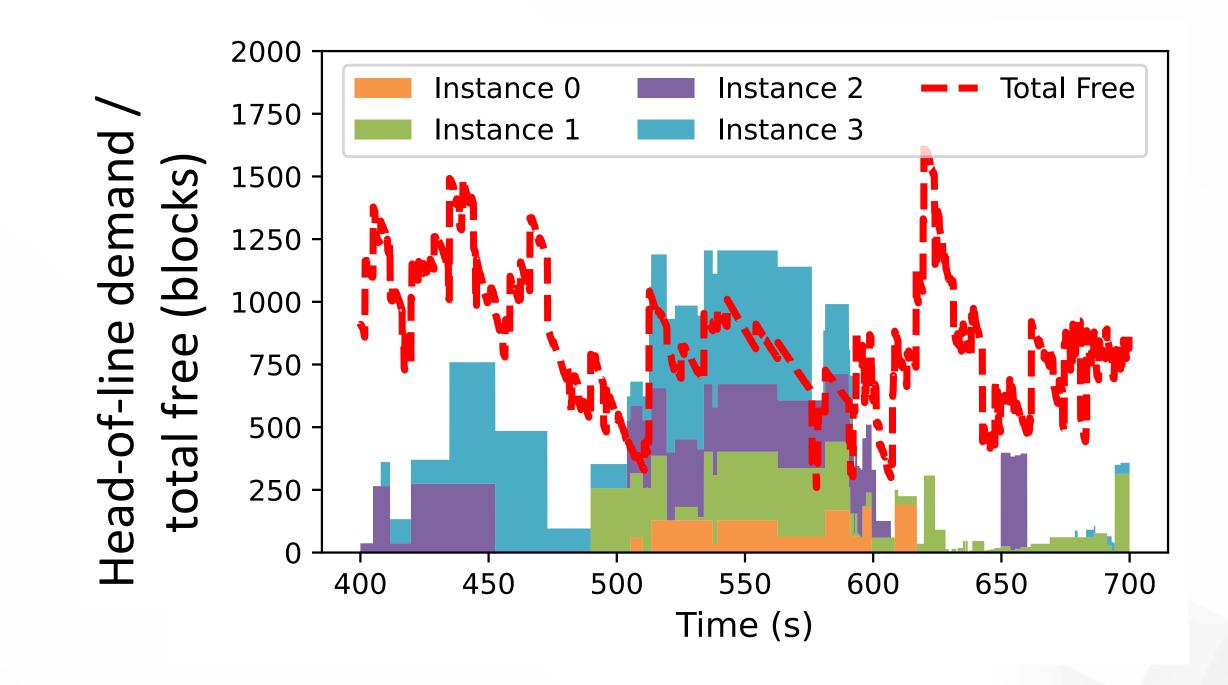


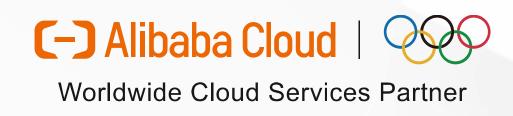
Packing



# **Challenge (2): Memory Fragmentation**

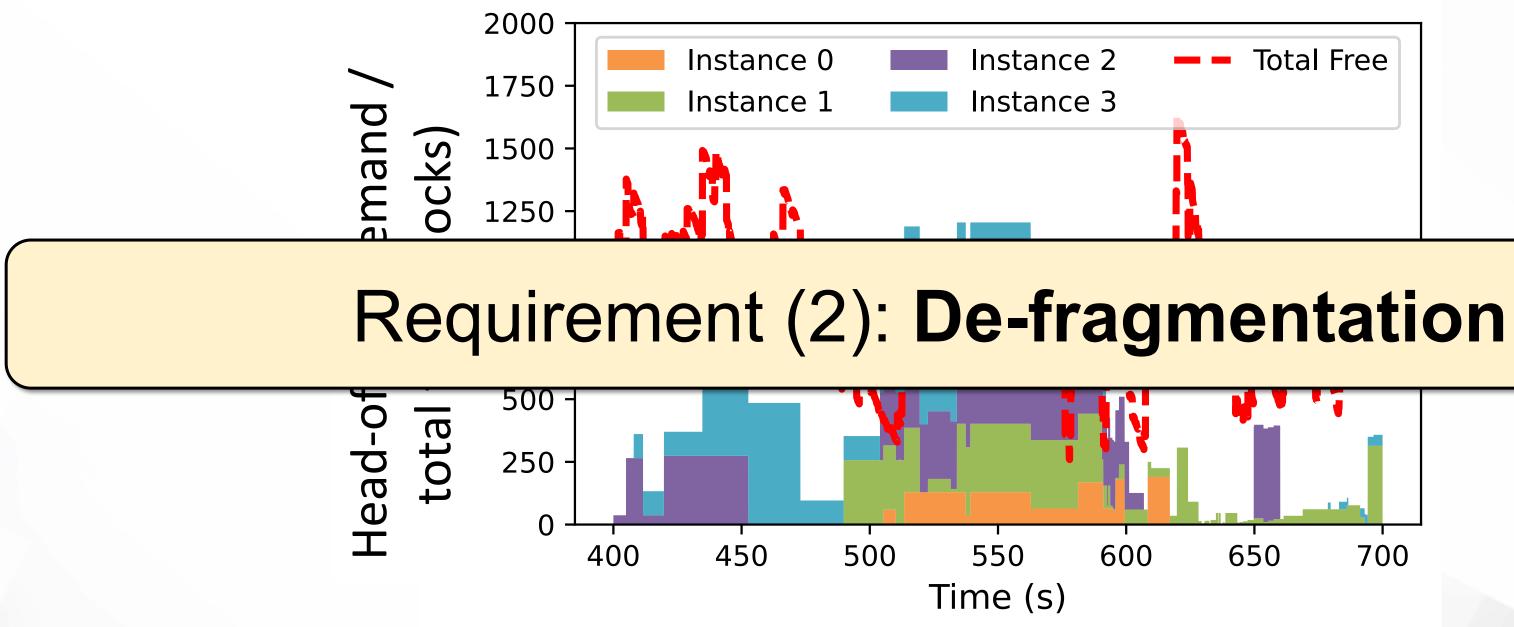
- Load balancing -> fragmentation across instances
  - A classic spreading vs. packing tradeoff
- Fragmentation -> worse queuing delays (first-token latencies)
  - A large space on one instance needed for the prompt

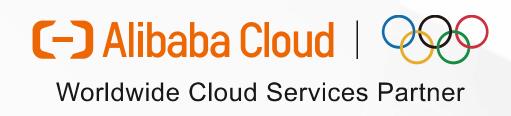




# **Challenge (2): Memory Fragmentation**

- Load balancing -> fragmentation across instances
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# **Challenge (3): Differentiated SLOs**

- Existing systems treat all requests equally
- Urgent requests could be easily interfered by normal ones
  - Queuing delays •
  - Performance interference

# Requirement (3): Request priorities



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# LLMs are *Multi-Tenant* and *Dynamic*



# A behavior that is:

# **Different from traditional DNNs**

- **Deterministic, stateless execution**
- Homogeneous requests

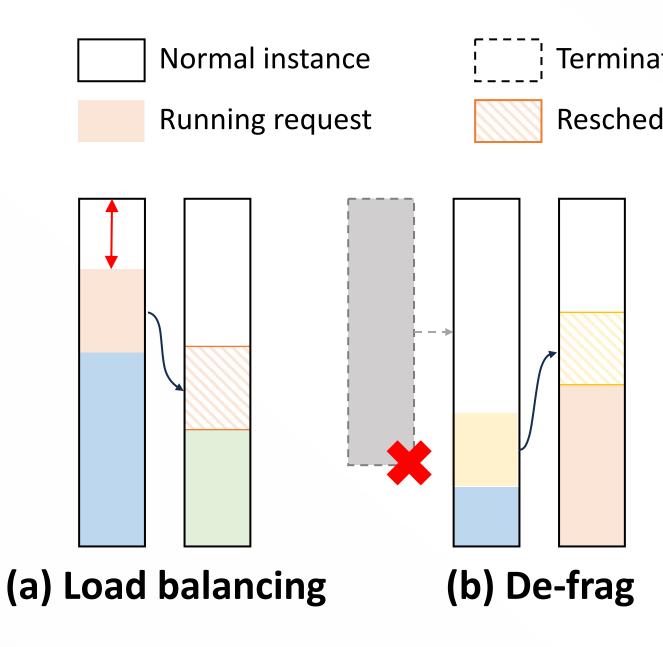
## but...

- Not new in modern operating / distributed systems Processes with dynamic working sets, unknown durations, different priorities, ...
- Context switching, process migration, ...



# Llumnix: Serving LLMs, the "OS" Way

- Continuous rescheduling across instances
  - Combined with dispatching and auto-scaling
- Powerful in various scheduling scenarios







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Terminating instance

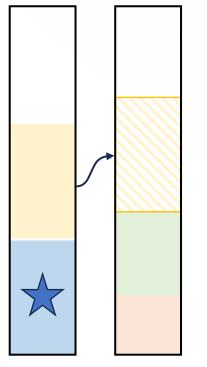


Queuing request

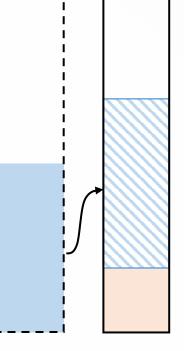
Rescheduling destination



High-priority request



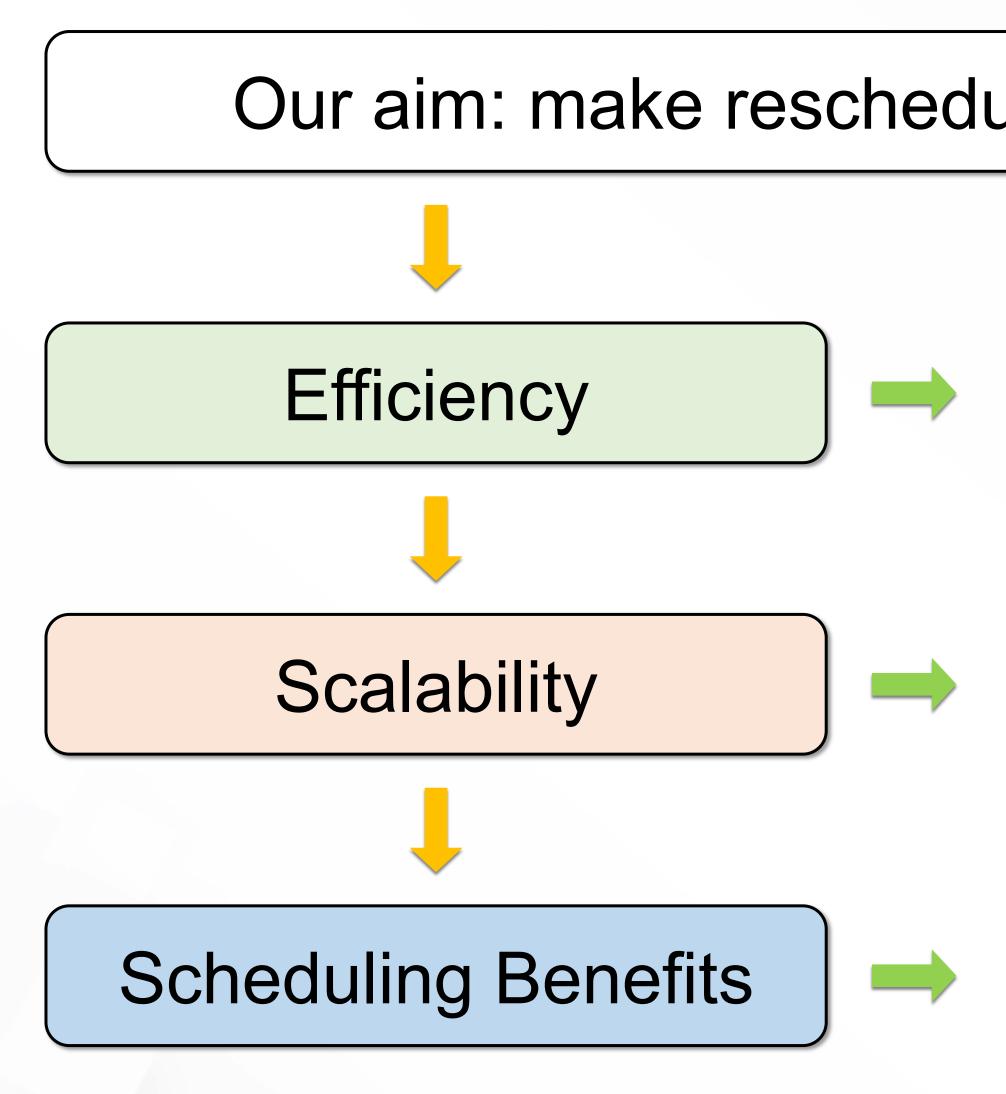
(c) Prioritization

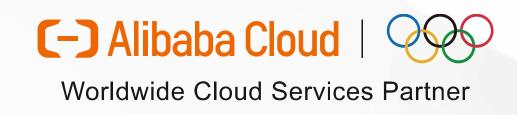


(d) Auto-scaling









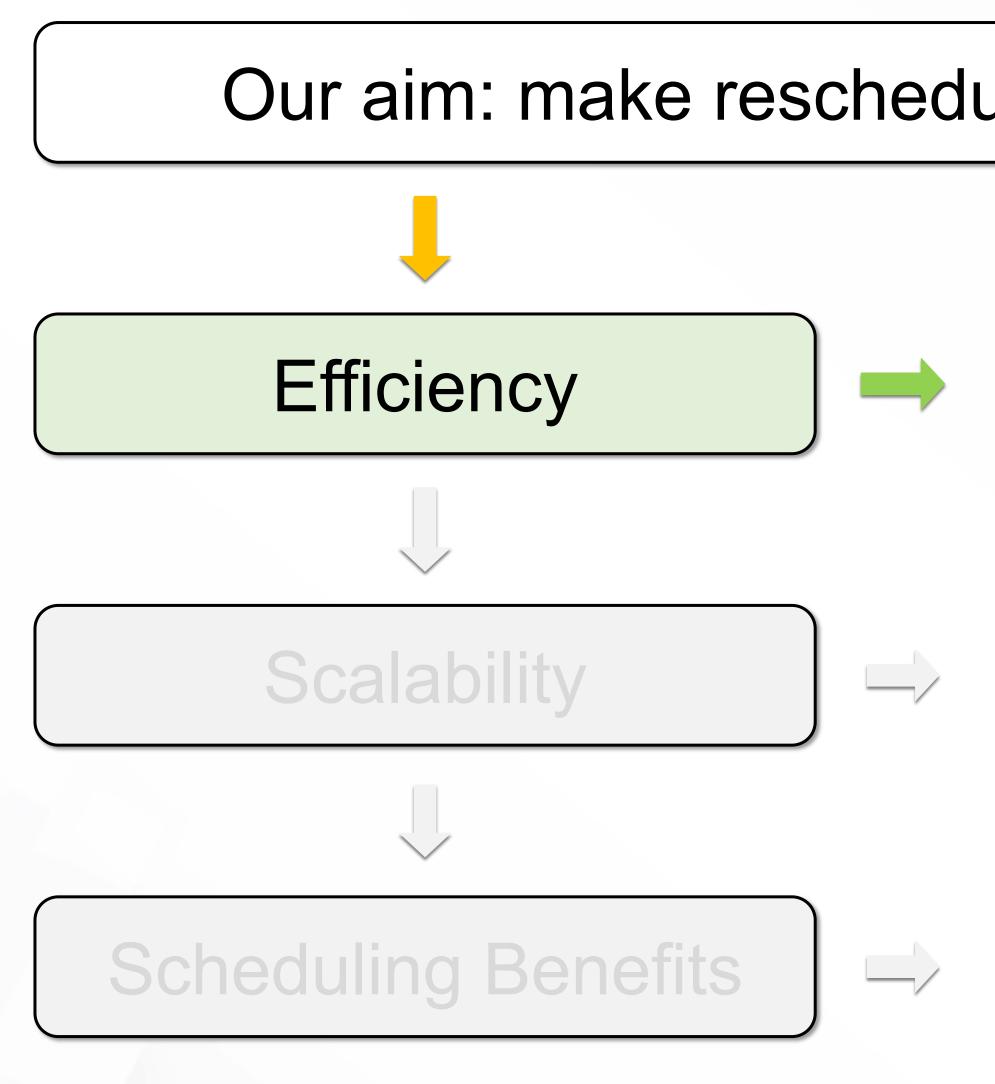
Our aim: make rescheduling the norm in LLM serving

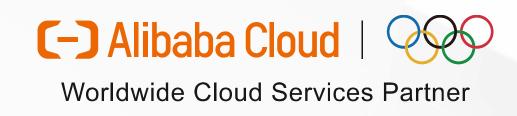
Live migration mechanism

## Distributed scheduling architecture

Unified, multi-objective scheduling policy







Our aim: make rescheduling the norm in LLM serving

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# **How to Reschedule KV Caches?**

# Overhead Performance

## Llumnix's live migration Near-zero downtime and overhead by design



**Rescheduling Downtime** 



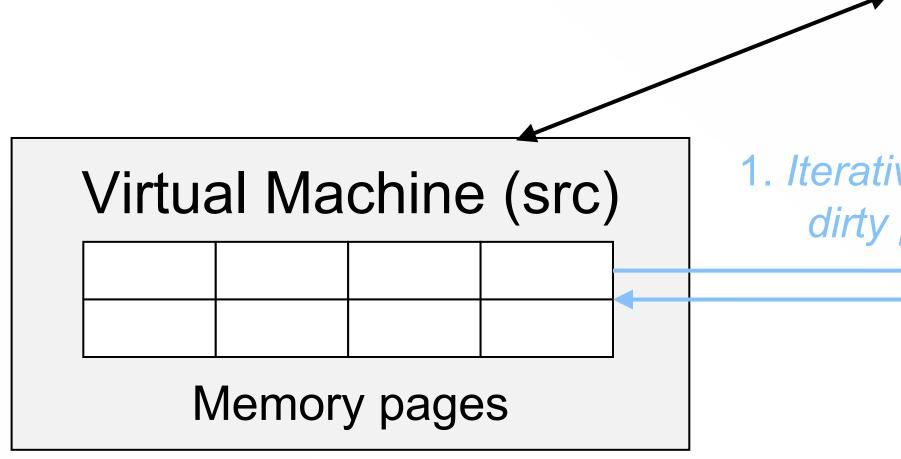
## Recompute

Downtime and overhead (compute waste) growing with sequence lengths

## Suspend-and-copy

Downtime for data transfer growing with sequence lengths





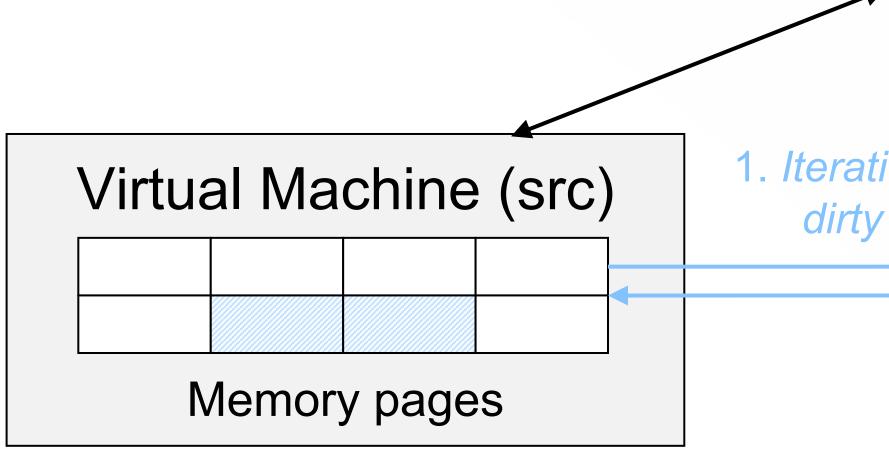




## 1. Iterative pre-copy with dirty page tracing

# Virtual Machine (dst)

## Memory pages



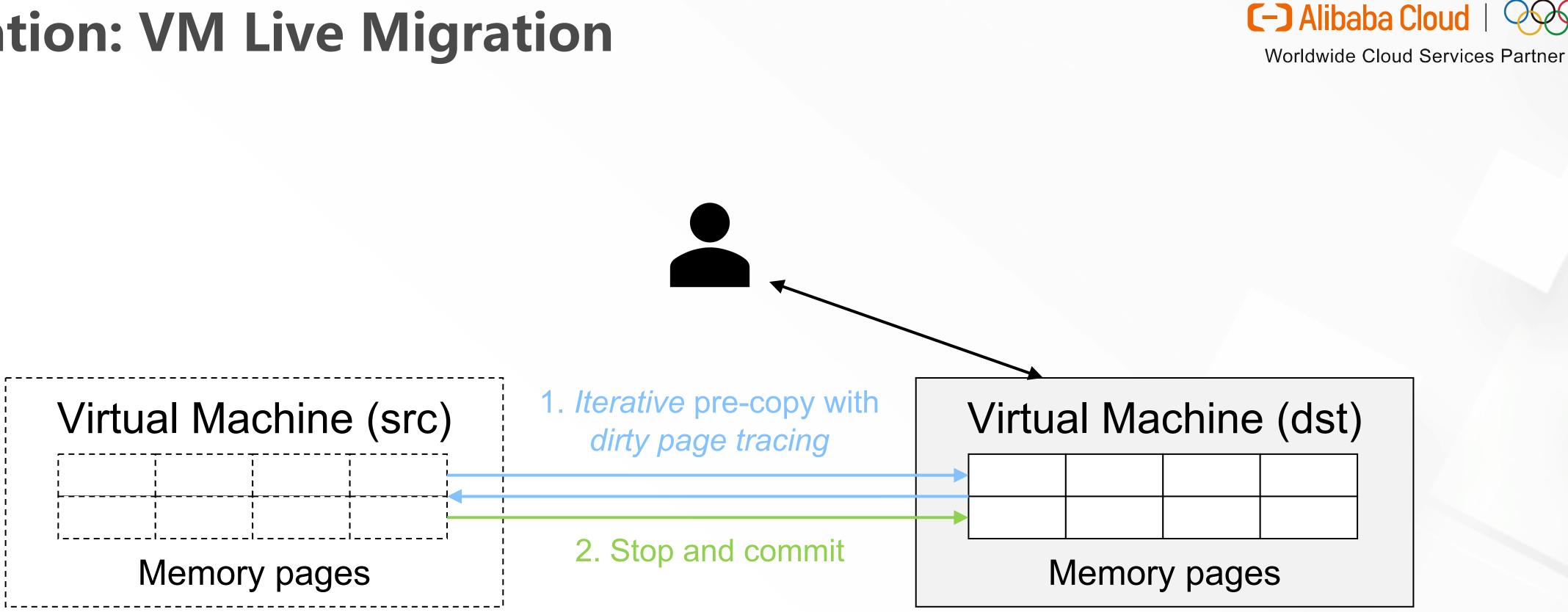




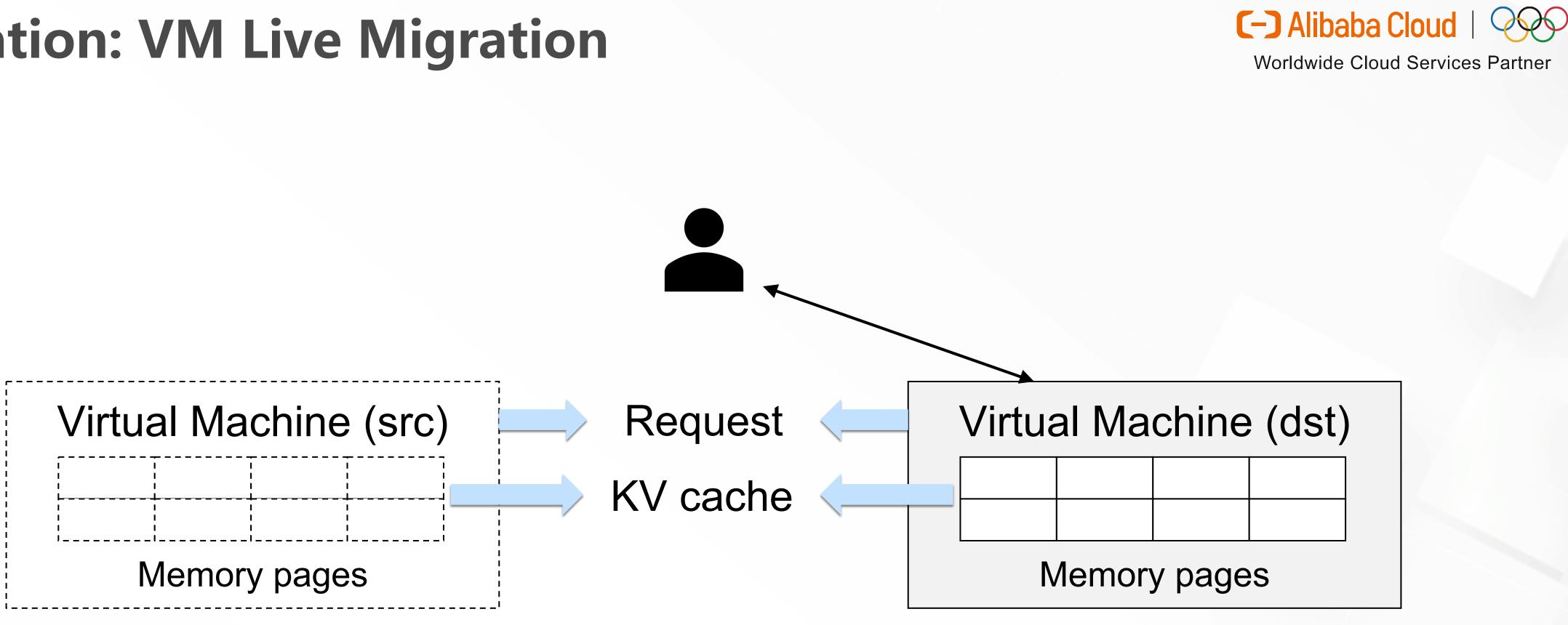
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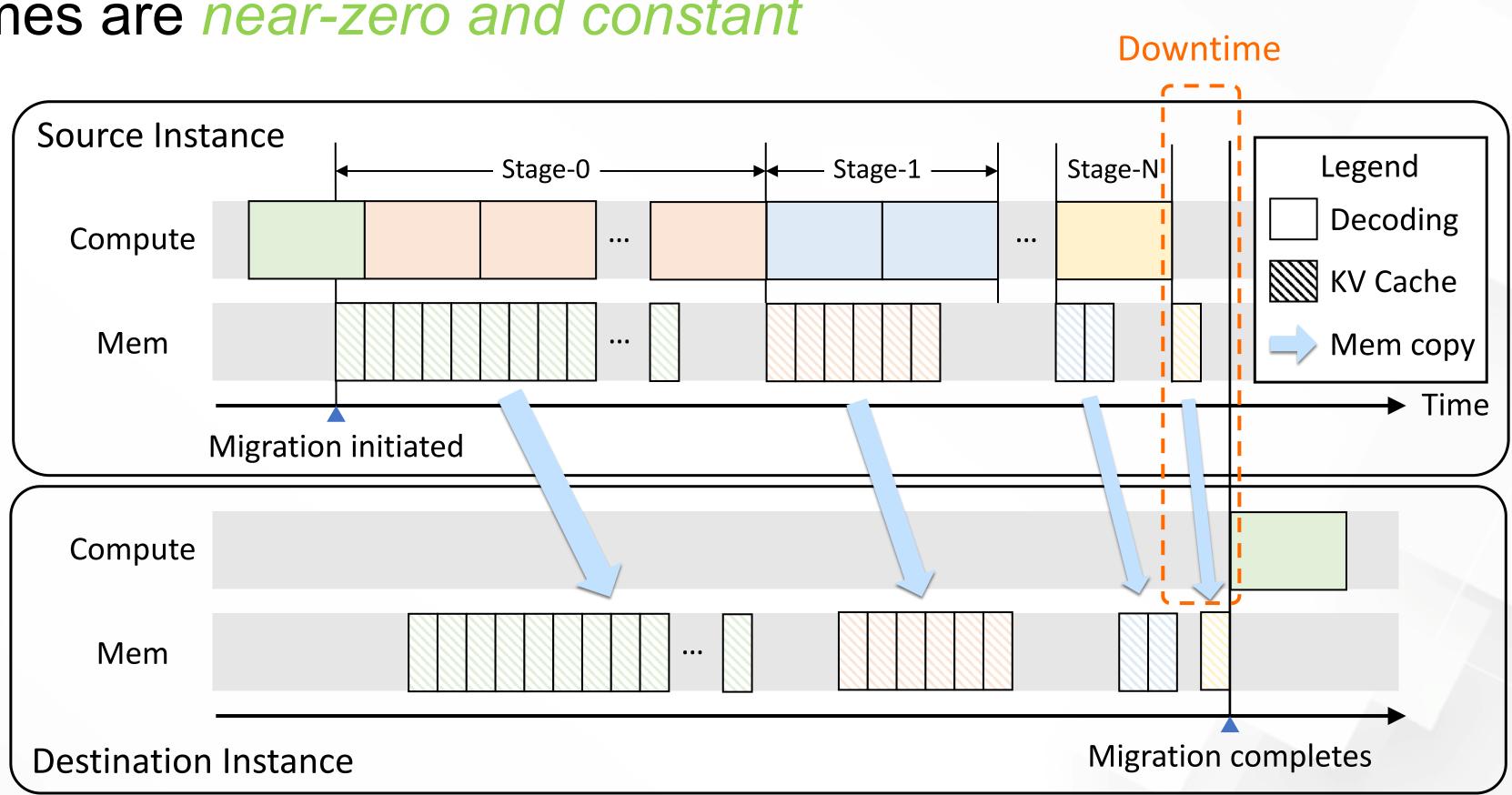




What are dirty pages?

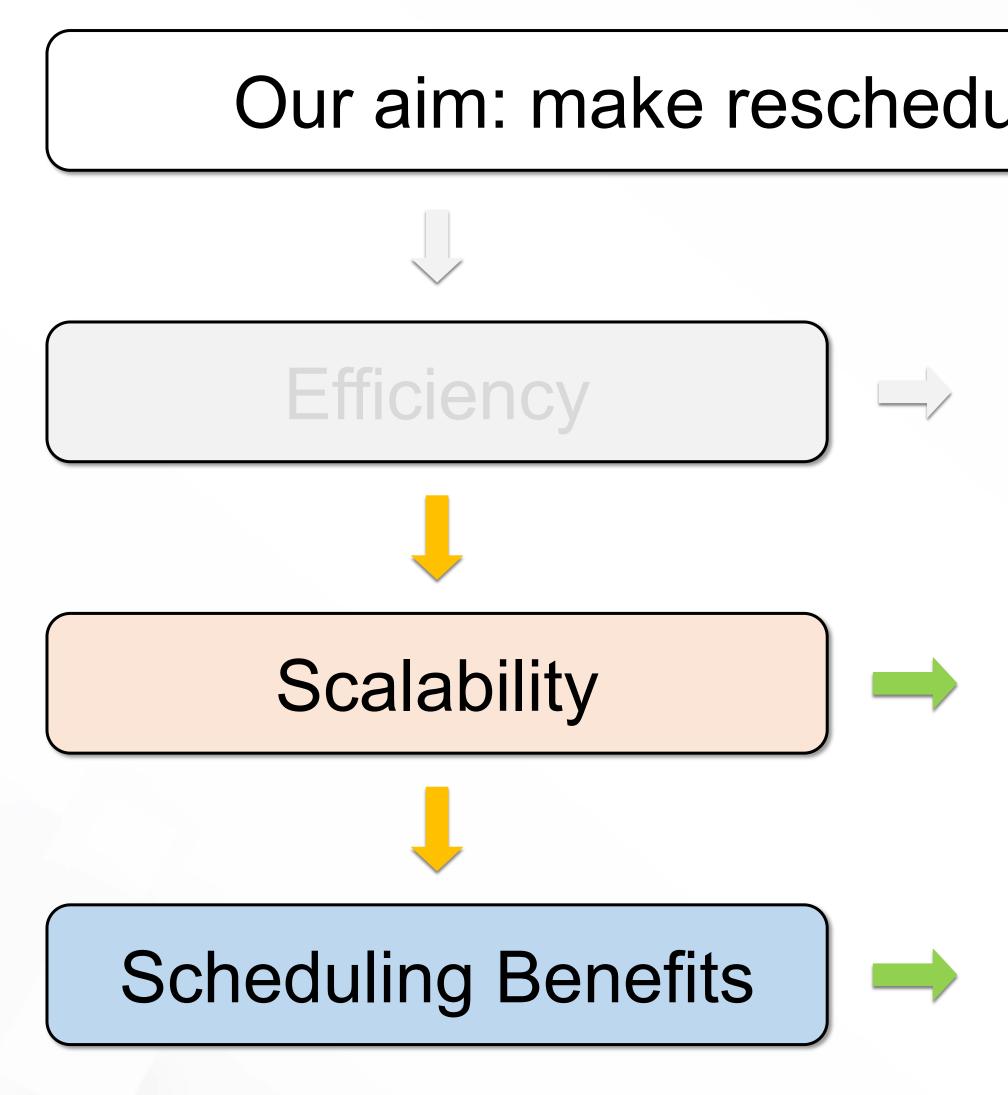
# Live Migration of LLM Requests

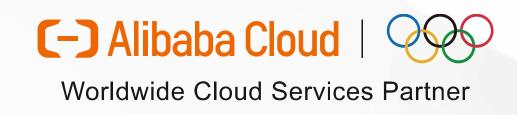
- KV caches are append-only
  - Copy dirty incremental blocks iteratively
  - Downtimes are *near-zero and constant* ullet











Our aim: make rescheduling the norm in LLM serving

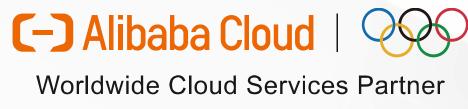
Live migration mechanism

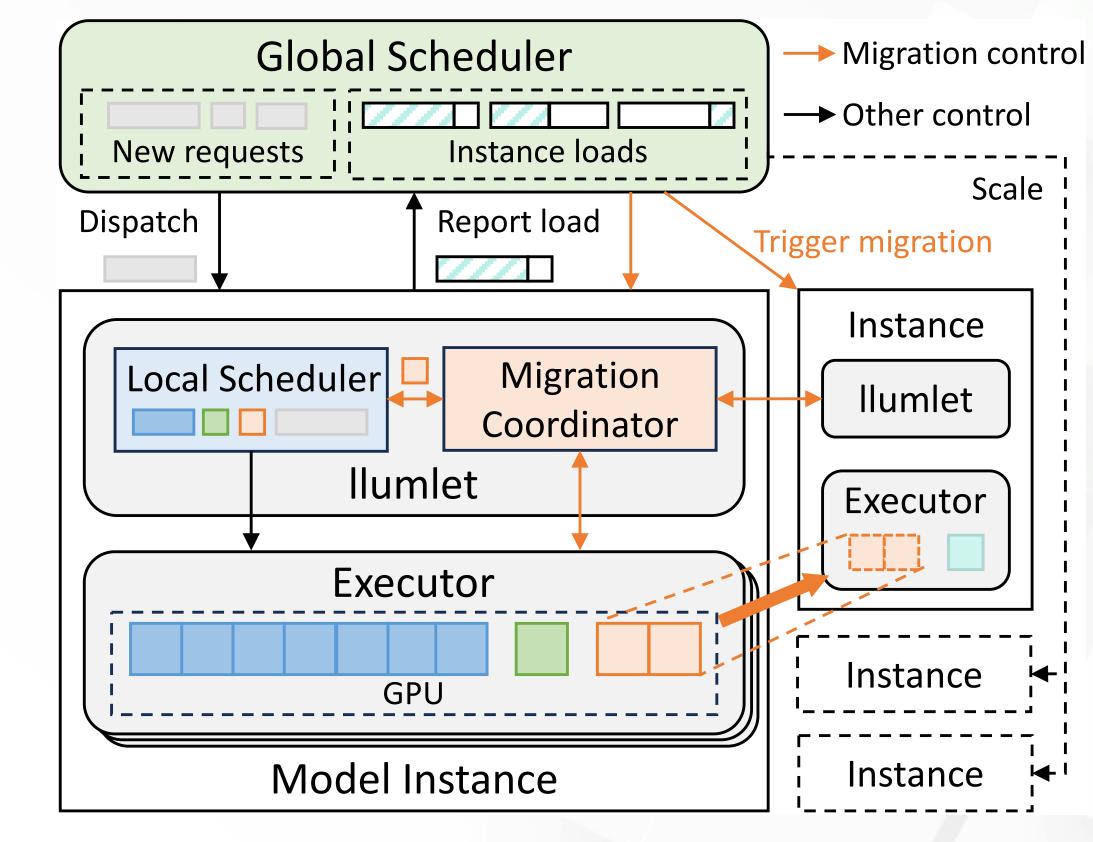
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# **Distributed Scheduling Architecture**

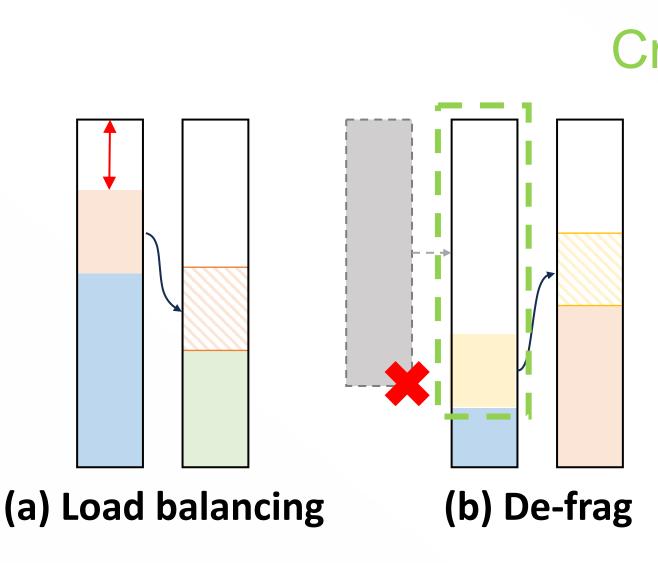
- Global scheduler for cross-instance scheduling
- Distributed **llumlets** for local scheduling
- A narrow interface: instance load





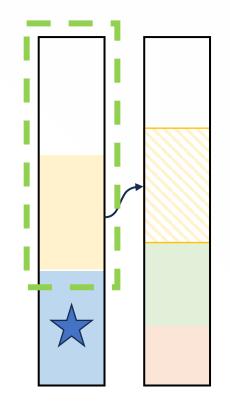
# **Scheduling Policy (sketch)**

- Virtual usage: unifying multiple objectives
- Policy: load balancing based on virtual usages

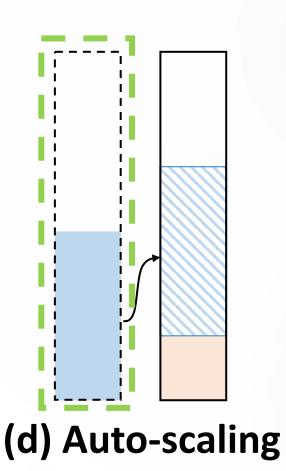




## Create free space

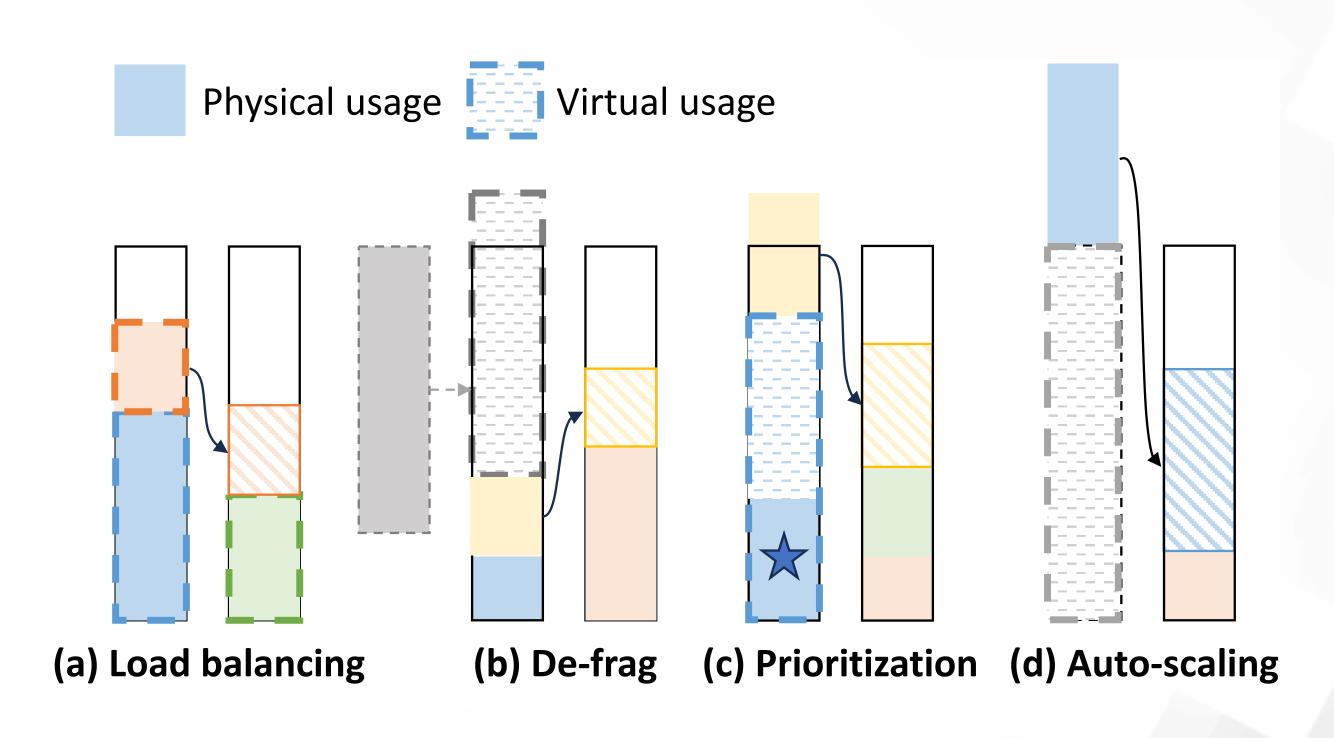


(c) Prioritization



# **Scheduling Policy (sketch)**

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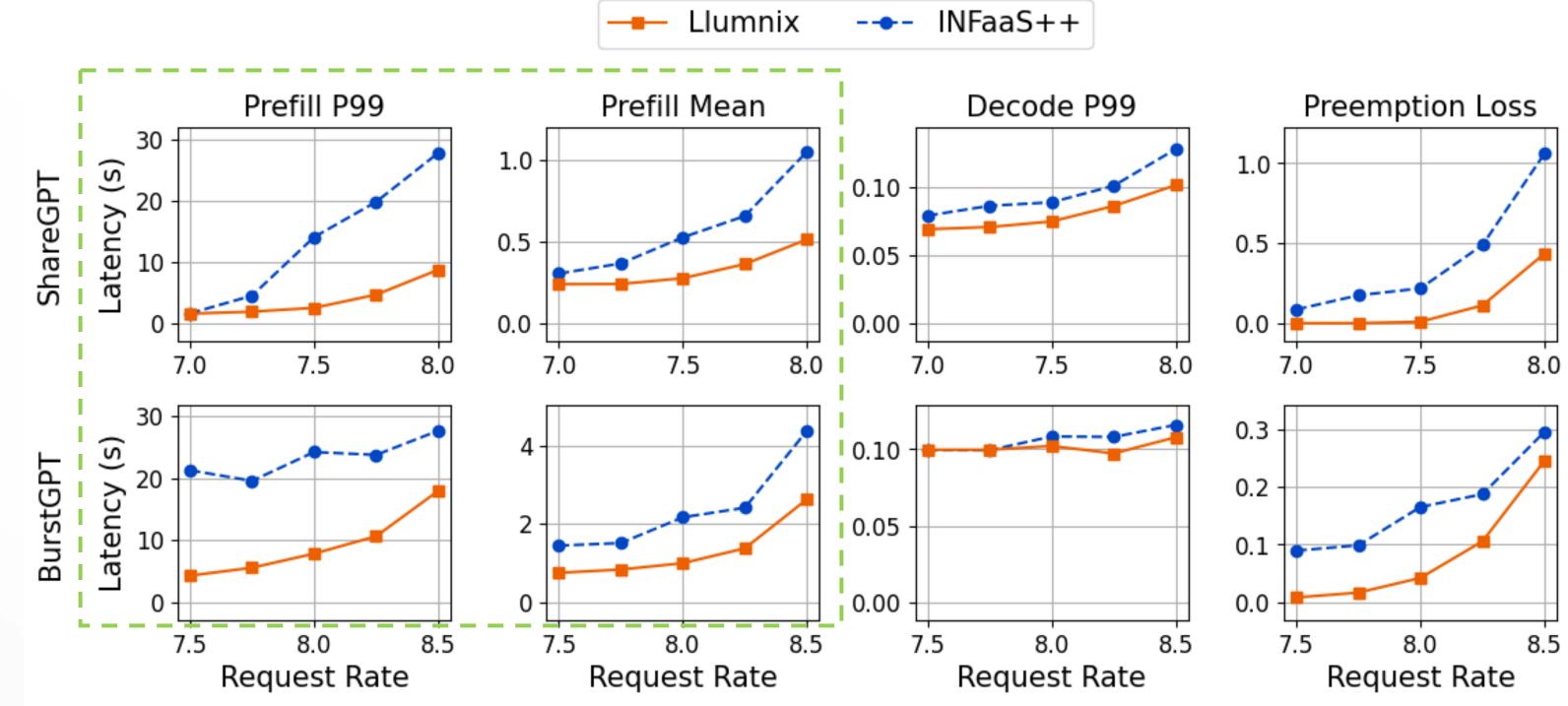
# **Evaluation**

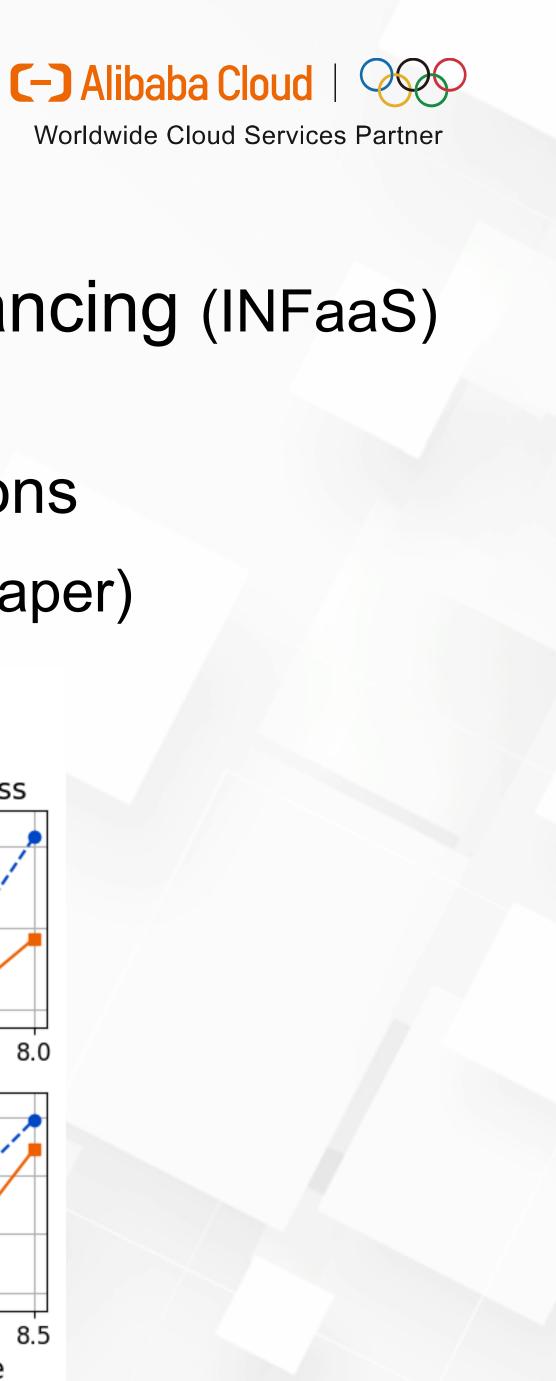
- Implemented as a scheduling layer atop vLLM
- Testbed: 16 A10 GPUs (24GB)
  - 4 4-GPU VMs, PCIe 4.0 in each node, 64Gb/s Ethernet across nodes
- Models: LLaMA-7B and LLaMA-30B
- Traces: ShareGPT, BurstGPT, generated power-law distributions



# End-to-end Serving Performance (16 LLaMA-7B instances)

- - Up to 2.2x/5.5x for first-token (mean/P99) via de-fragmentation
  - Up to 1.3x for per-token generation P99 via reducing preemptions
- More gains with more diverse sequence lengths (details in paper)

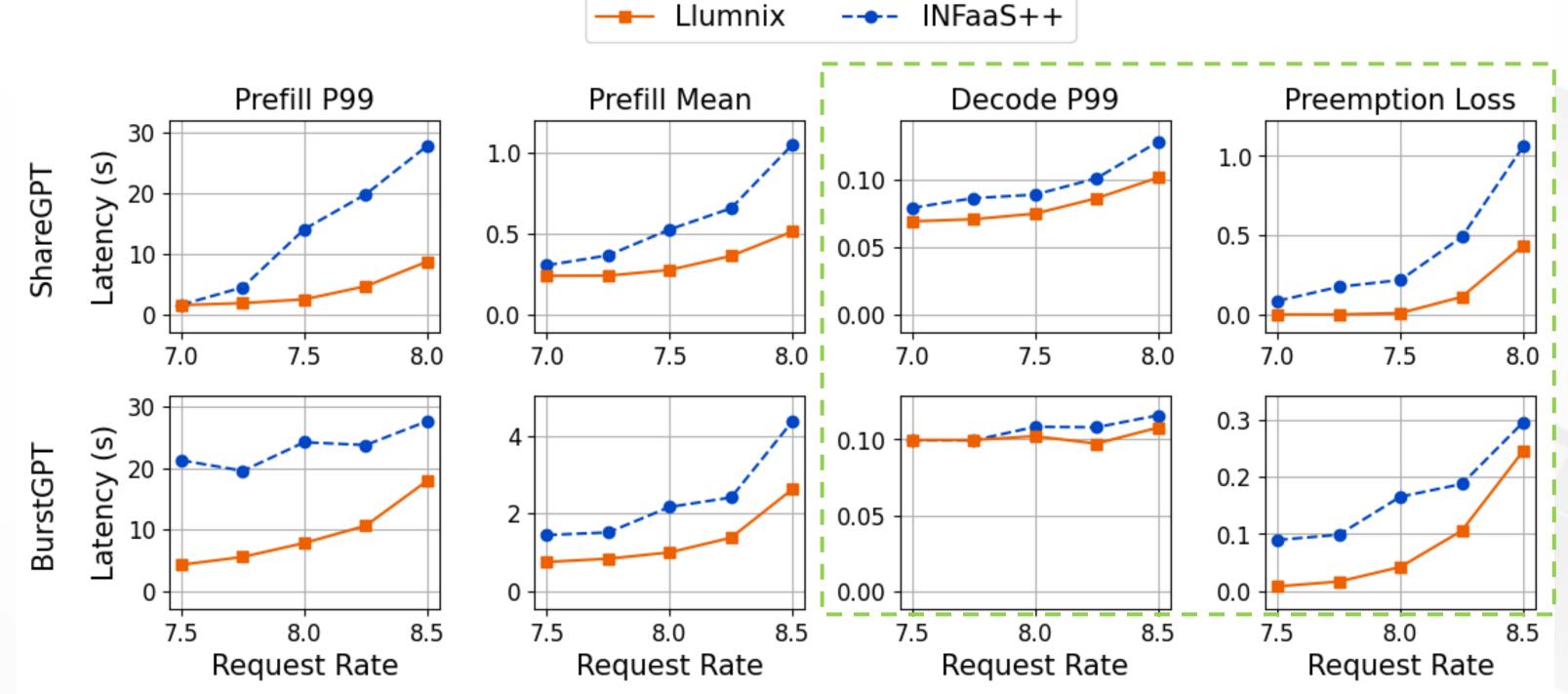


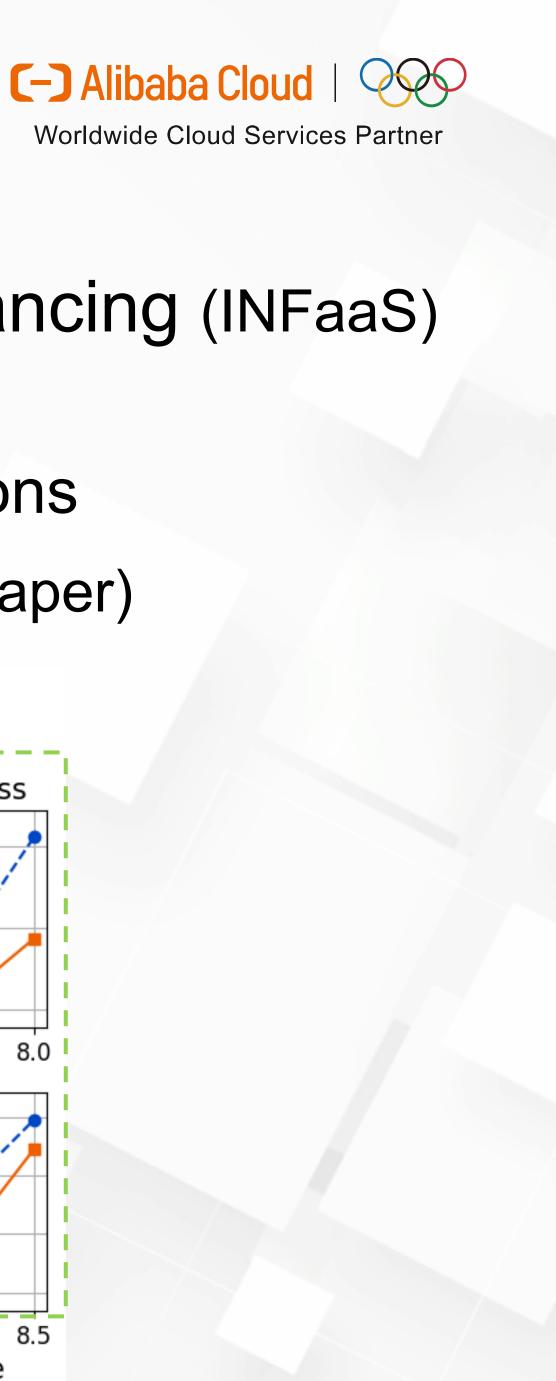


Benefits of migration: compared to dispatch-time load balancing (INFaaS)

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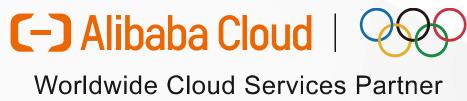


Benefits of migration: compared to dispatch-time load balancing (INFaaS)

# Conclusion

- Dynamic workloads need dynamic scheduling
  - LLMs are no exception
- Llumnix draws lessons from conventional systems wisdom
  - Classic scheduling goals in the new context of LLM serving •
  - Implementation of rescheduling with request live migration lacksquare
  - Continuous, dynamic rescheduling exploiting the migration ullet

https://github.com/AlibabaPAI/Ilumnix







# Thanks